# Concurrent access to shared mutable objects in Encore and its GC related implications

## 2015-09-28 Tobias Wrigstad && Albert Yang

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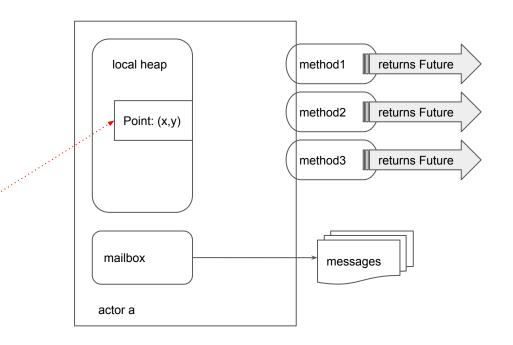
## Outline

- Encore Programming Model
- Sharing mutable objects
- How Object GC works in Encore
- Example of Object GC failure
- How Object GC in other actor-based systems work
- Shared Object
- Summary

#### **Encore Programming Model**

- Active objects (actor)
  - own thread of control
  - async interface
  - local heap
  - mailbox
- Passive objects (object)
  - like plain old Java object
  - sync interface
  - reside in local heap of actors

object



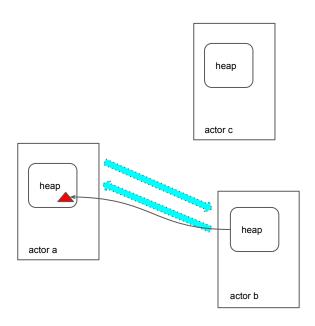
## Sharing mutable objects

- Using active objects
  - high latency if many actors scheduled in front
  - memory overhead; each object carries a local heap
  - serial access even for non-overlapping operations
- Using passive objects
  - Data Race
    - i. mutates objects after sending
    - ii. mutates objects after receiving
  - Breaks parallel GC
    - i. parallel GC on local heap of each actor
    - ii. PonyRT deals with sharing immutable objects
    - iii. but sharing mutable objects would lead to dangling pointers

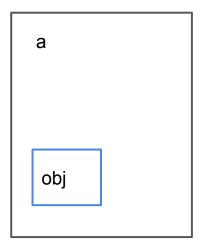
solvable using lock, STM, type system...

#### GC in Encore

- Actor GC [Clebsch, Drossopoulou '13]
  - detect unreachable actors and collect them
- Object GC [Clebsch et\_al. '15]
  - based on ownership and reference counting
  - local GC on the heap to each actor
  - sync using message passing



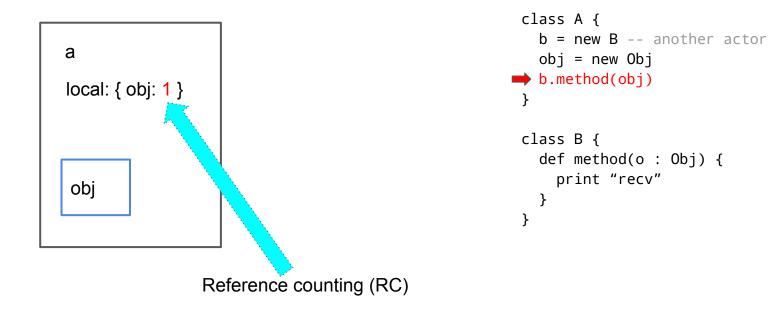
# Object GC (1)



```
class A {
   b = new B -- another actor
   obj = new Obj
   b.method(obj)
   }
class B {
   def method(o : Obj) {
     print "recv"
   }
  }
}
```

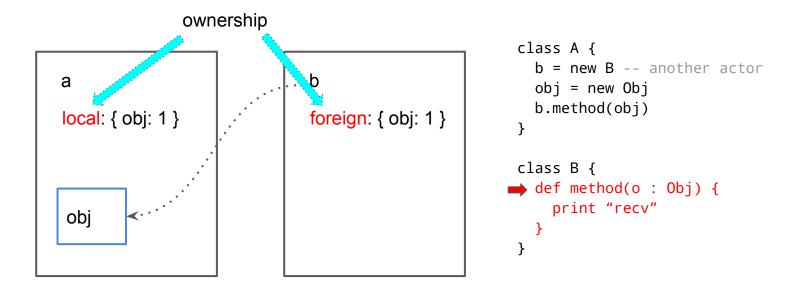
actor **a** creates **obj** in its local heap

# Object GC (2)



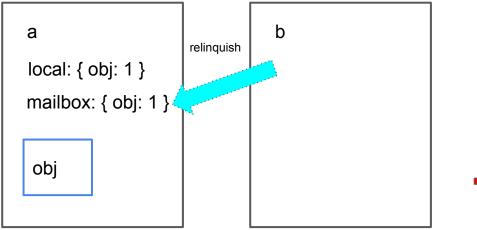
actor a creates an entry in **local** set before sending it out





actor b receives obj and creates an entry in foreign set, for it doesn't belong to the current actor

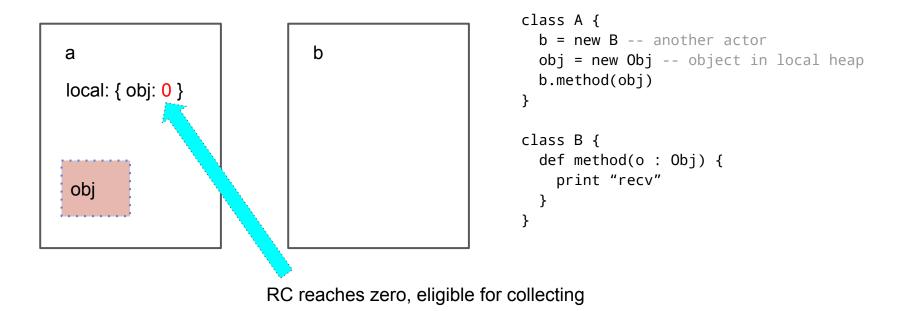
# Object GC (4)



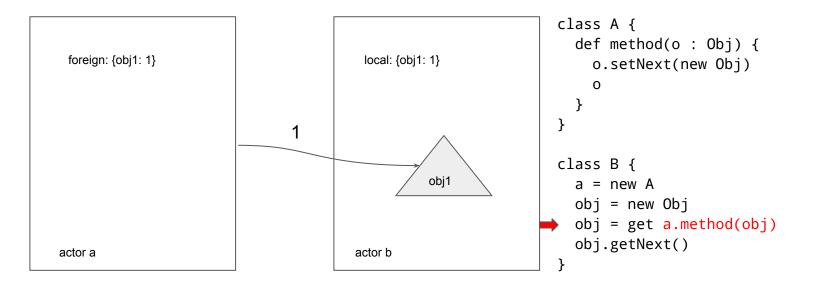
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   b = new B -- another actor
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   def method(o : Obj) {
      print "recv"
   }
}
```

During GC in **b**, **obj** is found unreachable and reports to the owner, **a** 

# Object GC (5)

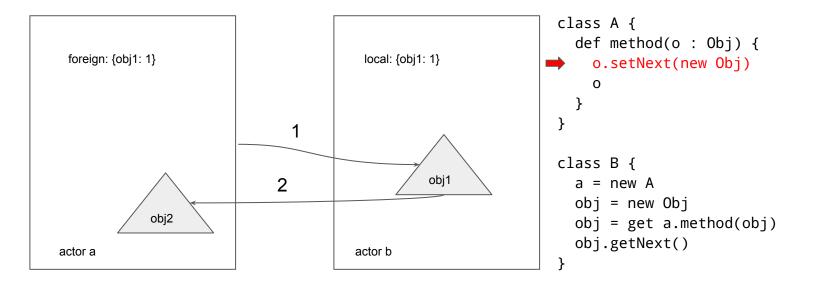


## Example of extending objects after receiving (1)

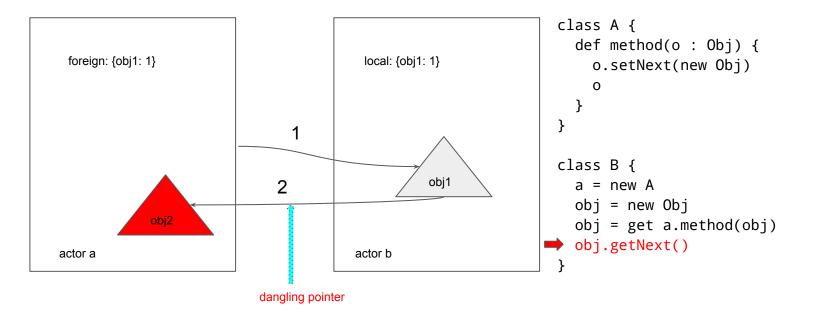


actor b creates obj1 and shares with actor a

## Example of extending objects after receiving (2)

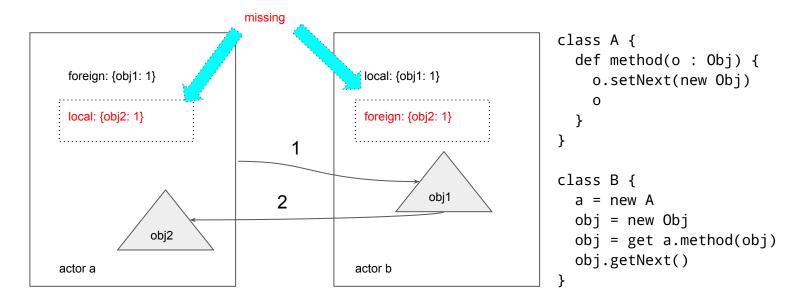


#### Example of extending objects after receiving (3)



GC in actor a would collect obj2, because it's not captured in the local set

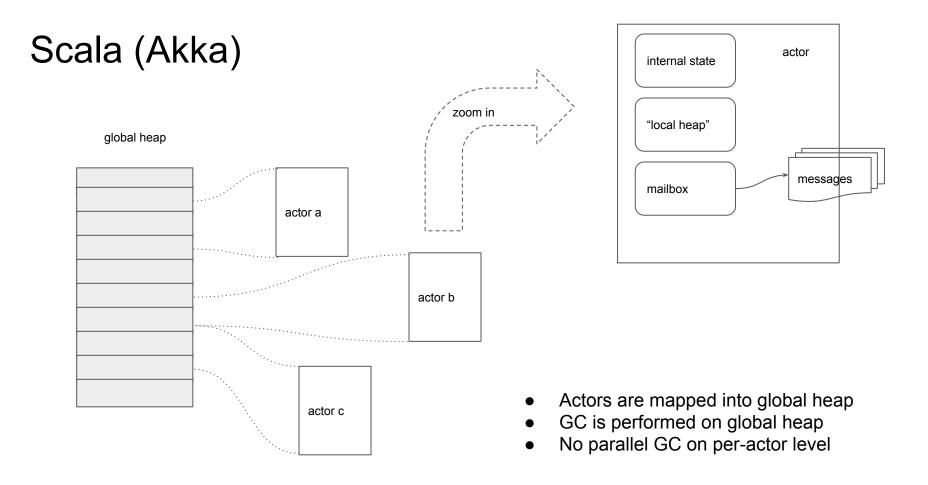
#### Example of extending objects after receiving (4)



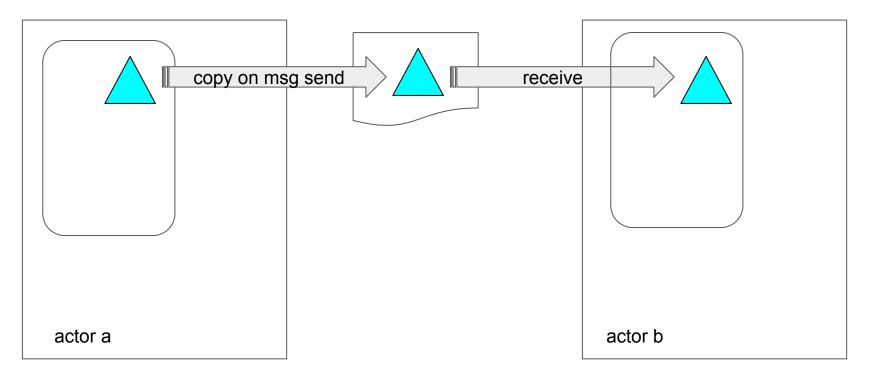
The correct state to keep obj2 from being collected

#### How Object GC in other actor-based systems work

- Scala (Akka)
- Erlang



# Erlang (1)



# Erlang (2)



- Two independent copies
- Parallel GC without interference

## Comparison on Object GC

- Scala (Akka)
  - global GC can't take advantage of the isolation of actors
- Erlang
  - local GC, but sharing large objects is expensive
- Encore
  - local GC, sharing is cheap, but has GC issues in concurrent access to shared passive objects...

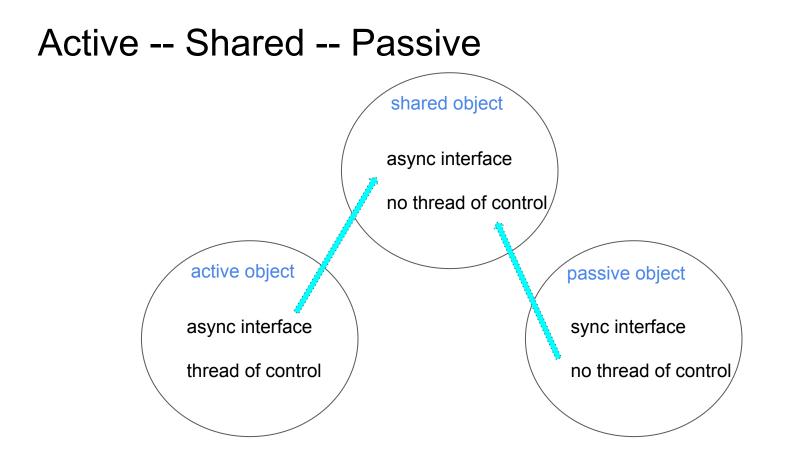
#### Proposed Solution: Shared Object

Isolated container, supporting a limited form of sharing and contention

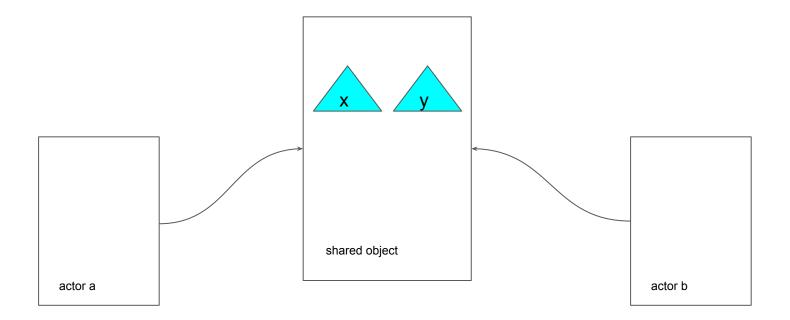
- Support concurrent access to **its** encapsulated passive objects
- Support for parallel method execution

Important GC points

- No global sharing of mutable state, i.e., still parallel GC
- Only place where GC is contended is inside the heaps of shared objects

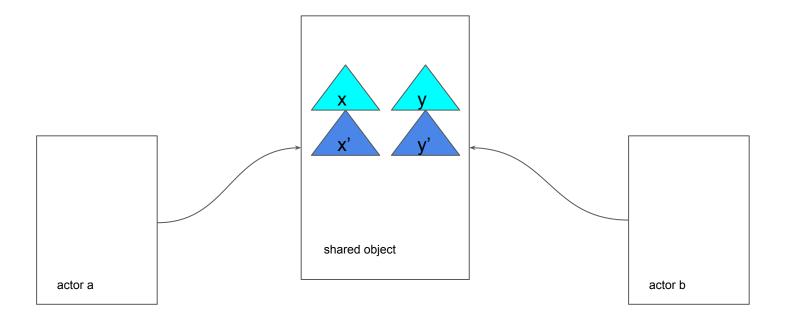


Example of Shared Object (1)



actor **a** and **b** have access to a shared object, that has two disjoint fields, **x** and **y** 

Example of Shared Object (2)



two actors could extend mutable state on shared object's heap, synchronously and concurrently

#### **Implementation Space**

- Actor
  - Shared Object has the same interface as actor
  - high latency due to async operations
- Semi-actor
  - the caller actor acquires the Shared Object if it's not acquired already
  - the caller actor do operations synchronously
  - the caller actor insert the operations to Shared Object's mailbox it's already acquired
  - (Similar to Queue Delegation locking.)
- STM
  - Wrap each operation inside a transaction
  - get parallelism when operations don't overlap

## Summary

- GC trade-offs: local vs global heap, copy vs pass-by-ref
- Sharing mutable passive objects could cause GC issues in local + pass-by-ref
- Shared Objects (SO) enables shared states in its heap
- Actors could collaborate on data structures in parallel via SO
- Large implementation space: actor, QD locking, STM...

Questions, Comments Suggestions