

# The Rödl Nibble

Frédéric Havet

COATI, INRIA, I3S, CNRS, Univ. Nice Sophia Antipolis  
Sophia Antipolis, France

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# Beyond design

Design = perfect case

**BUT**

do not always exist.

No decomposition of  $K_5$  into copies of  $K_3$ .

$$(|E(K_5)| = 10 \text{ and } |E(K_3)| = 3.)$$

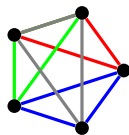
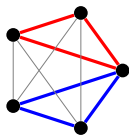
No decomposition of  $K_6^{(3)}$  into copies of  $K_4^{(3)}$ .

$$(\text{degree in } K_6^{(3)} = 10 \text{ and degree in } K_4^{(3)} = 3)$$

# Approximate design

**packing** of  $F$  in  $G$  : set of edge-disjoint copies of  $F$  in  $G$ .

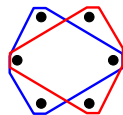
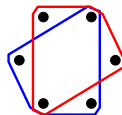
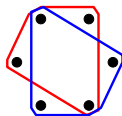
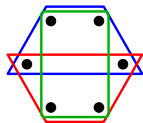
**covering** of  $G$  by  $F$  : set of copies of  $F$  such that every edge of  $G$  is in one of the copies.



# Approximate design

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# Approximate design

**packing** of  $F$  in  $G$  : set of edge-disjoint copies of  $F$  in  $G$ .

**covering** of  $G$  by  $F$  : set of copies of  $F$  such that every edge of  $G$  is in one of the copies.

$m(n, k, t)$  : maximum size of a packing of  $K_k^{(t)}$  in  $K_n^{(t)}$

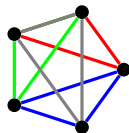
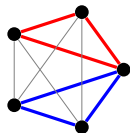
$M(n, k, t)$  : minimum size of a covering of  $K_n^{(t)}$  by  $K_k^{(t)}$ .

# Trivial inequalities

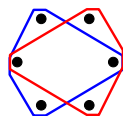
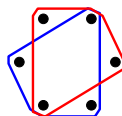
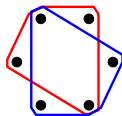
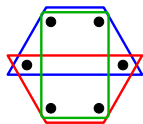
$$m(n, k, t) \leq \frac{\binom{n}{t}}{\binom{k}{t}} \leq M(n, k, t)$$

with equalities iff there is a  $t$ - $(n, k, 1)$ -design.

# Examples



$$m(5, 3, 2) = 2 < \frac{\binom{5}{2}}{\binom{3}{2}} = \frac{10}{3} \leq M(5, 3, 2) = 4$$



$$m(6, 4, 3) = 3 < \frac{\binom{6}{3}}{\binom{4}{3}} = 5 < M(6, 4, 3) = 6$$

# Erdős–Hanani Conjecture

$$m(n, k, t) \leq \frac{\binom{n}{t}}{\binom{k}{t}} \leq M(n, k, t)$$

Conjecture (Erdős–Hanani, 1963)

$$\lim_{n \rightarrow +\infty} \frac{m(n, k, t)}{\binom{n}{t} / \binom{k}{t}} = 1 \quad \text{and} \quad \lim_{n \rightarrow +\infty} \frac{M(n, k, t)}{\binom{n}{t} / \binom{k}{t}} = 1.$$

Erdős–Hanani: True for  $t = 2$  (graphs).



# Rödl Theorem

$$m(n, k, t) \leq \frac{\binom{n}{t}}{\binom{k}{t}} \leq M(n, k, t)$$

## Theorem (Rödl, 1985)

$$\lim_{n \rightarrow +\infty} \frac{m(n, k, t)}{\binom{n}{t} / \binom{k}{t}} = 1 \quad \text{and} \quad \lim_{n \rightarrow +\infty} \frac{M(n, k, t)}{\binom{n}{t} / \binom{k}{t}} = 1.$$

# The auxiliary hypergraph $\mathcal{H}$

**Hypergraph  $\mathcal{H}$**  : vertices =  $t$ -subsets of  $[n]$   
hyperedges = sets of  $\binom{k}{t}$   $t$ -subsets of a  $k$ -subset of  $[n]$ .

packing of  $K_k^{(t)}$  in  $K_n^{(t)}$   $\leftrightarrow$  **matching** in  $\mathcal{H}$

**matching** : set of pairwise vertex-disjoint hyperedges.

covering of  $K_n^{(t)}$  by  $K_k^{(t)}$   $\leftrightarrow$  **cover** of  $\mathcal{H}$

**cover** : set of hyperedges s.t. every vertex is in one of them.

## Theorem (Rödl, 1985)

$\mathcal{H}$  has a matching of size  $(1 - o(1)) N/r$   
and a cover of size  $(1 - o(1)) N/r$

$$N = \binom{n}{t}$$
$$r = \binom{k}{t}.$$

## A more general result

- $\mathcal{H}$  has  $N$  vertices with  $N = \binom{n}{t}$ .
- $\mathcal{H}$  is  $r$ -uniform with  $r = \binom{k}{t}$ .
- $\mathcal{H}$  is  $D$ -regular with  $D = \binom{n-t}{k-t}$ .

**Observation** : The existence of the desired matching and cover also holds for large class of  $r$ -uniform hypergraphs  $D$ -regular graphs.

Frank and Rödl, Pippenger and Spencer 1989, Kahn 1996.

# Equivalence of the two statements

In a  $D$ -regular  $r$ -uniform hypergraph of order  $N$

matching of size  $(1 - o(1)) N/r \Leftrightarrow$  cover of size  $(1 + o(1)) N/r$

$(\Rightarrow)$  matching of size  $\frac{N(1-\epsilon)}{r}$ ,  $\Rightarrow$  at most  $\epsilon N$  non-covered vertices  
 $\Rightarrow$  cover of size at most  $\frac{N(1-\epsilon)}{r} + \epsilon N$ .

$(\Leftarrow)$  cover of size  $\frac{N(1+\epsilon)}{r}$ . Each vertex  $x$  covered  $c(x)$  times.

For each  $x$  choose  $c(x) - 1$  hyperedges and remove them.

At most  $\sum_{v \in V} (c(x) - 1) = \frac{(1+\epsilon)N}{r} \times r - N = \epsilon N$  are removed.

$\Rightarrow$  matching of size at least  $\frac{N(1+\epsilon)}{r} - \epsilon N$ .

## General idea: nibbling

Fix  $\epsilon$ .

Take a random set of  $\epsilon N/r$  edges. W.h.p. only  $O(\epsilon^2 N)$  vertices covered more than once. So at least  $\epsilon N - O(\epsilon^2 N)$  covered vertices.

Remove the covered vertices.

Choose again a random set of edges covering roughly an  $\epsilon$ -fraction of the vertices with almost no overlap.

And so on until at most  $\epsilon N$  vertices remain.

Cover each of the remaining vertices with a dedicated hyperedge.

# The Theorem

$r > 2$  fixed.

For  $\kappa \geq 1$  and  $a > 0$ , there exists  $\gamma > 0$  and  $d_0$  s.t. the following holds for all  $N \geq D \geq d_0$ .

Every  **$r$ -uniform hypergraph**  $H = (V, \mathcal{E})$  on  $N$  vertices such that

- (0)  $d(x) > 0$  for all  $x \in V$ .
- (1) All vertices  $x$  except at most  $\gamma N$  satisfy  $d(x) = (1 \pm \gamma)D$ .
- (2) For all  $x \in V$ ,  $d(x) < \kappa D$ .
- (3) For each pair  $(x, y)$  of distinct vertices,  $d(x, y) < \gamma D$ .

has a **cover of size**  $(1 + a)\frac{N}{r}$ .

$d(x, y)$  : **codegree** = number of edges containing both  $x$  and  $y$ .

In  $\mathcal{H}$ , we have  $d(x, y) = \binom{n-t-1}{k-t-1} = o(D)$ .

# The Nibble Lemma

$r > 2$  fixed.

For  $K \geq 1$ ,  $\epsilon > 0$  and  $\delta' > 0$ , there exists  $\delta > 0$  and  $D^*$  s.t. the following holds for all  $N \geq D \geq D^*$ .

Every  $r$ -uniform hypergraph  $H = (V, \mathcal{E})$  on  $N$  vertices such that

- (1) All vertices  $x$  except at most  $\delta N$  satisfy  $d(x) = (1 \pm \delta)D$ .
- (2) For all  $x \in V$ ,  $d(x) < KD$ .
- (3) For each pair  $(x, y)$  of distinct vertices,  $d(x, y) < \delta D$ .

contains a set  $\mathcal{E}'$  of hyperedges s.t.

- (iv)  $|\mathcal{E}'| = \frac{N}{r}\epsilon(1 \pm \delta')$ .
- (v)  $V' = V \setminus \bigcup_{S \in \mathcal{E}'} S$  has size  $Ne^{-\epsilon}(1 \pm \delta')$ .
- (vi) All vertices  $x$  of  $V'$  except at most  $\delta'|V'|$  the degree  $d'(x)$  of  $x$  in  $H[V']$  satisfies  $d'(x) = De^{-\epsilon(r-1)}(1 \pm \delta')$ .

# Proving the Theorem with the Nibble Lemma

$a > 0$ . Take  $\delta, \epsilon$  very small, and  $p$  s.t.  $e^{-\epsilon p} < \epsilon$ .

$K_i = \kappa e^{i(r-1)}$ ,  $D_i = D e^{-\epsilon i(r-1)}$  with  $D \geq d_0$ .

Choose  $\delta = \delta_p > \delta_{p-1} > \dots > \delta_0$  s. t. we can apply the Nibble Lemma each time and  $\delta_{i-1} \leq \delta_i e^{-\epsilon(r-1)}$ .

Step  $i$ :  $H_{i-1} \rightsquigarrow H_i$ .

- (1) All  $x$  but at most  $\delta_{i-1} N_{i-1}$  satisfy  $d(x) = (1 \pm \delta_{i-1}) D_{i-1}$ .
- (2) For all  $x \in V_{i-1}$ ,  $d(x) < K_{i-1} D_{i-1}$ .
- (3) For each pair  $(x, y)$  of distinct vertices,  $d(x, y) < \delta_{i-1} D_{i-1}$ .

A set  $\mathcal{E}'_i$  of hyperedges s.t.

(iv)  $|\mathcal{E}'_i| = \frac{|V_{i-1}|}{r} \epsilon (1 \pm \delta_i)$ .

(v)  $V_i = V_{i-1} \setminus \bigcup_{S \in \mathcal{E}'_i} S$  has size  $|V_{i-1}| e^{-\epsilon} (1 \pm \delta_i)$ .

(vi) All vertices  $x$  of  $V_i$  except at most  $\delta_i |V_i|$  the degree of  $x$  in  $H_i$  satisfies  $d(x) = D_{i-1} e^{-\epsilon(r-1)} (1 \pm \delta_i) = (1 \pm \delta_i) D_i$ .



# Proving the Theorem with the Nibble Lemma

$a > 0$ . Take  $\delta, \epsilon$  very small, and  $p$  s.t.  $e^{-\epsilon p} < \epsilon$ .  
 $K_i = \kappa e^{i(r-1)}, \quad D_i = D e^{-i(r-1)}$  with  $D \geq d_0$ .

Choose  $\delta = \delta_p > \delta_{p-1} > \dots > \delta_0$  s. t. we can apply the Nibble Lemma each time and  $\delta_{i-1} \leq \delta_i e^{-\epsilon(r-1)}$ .

$$P = \prod_{i=0}^p (1 + \delta_i) \leq \prod_{i=0}^p (1 + \delta e^{-i\epsilon(r-1)}) \leq \frac{1+4\delta}{1+2\delta}.$$

$$|V_i| \leq N e^{-i\epsilon} P \leq N e^{-i\epsilon} (1 + 2\delta).$$

$$\begin{aligned} \text{So } |\mathcal{E}'_i| &= \left(\frac{\epsilon}{r} |V_{i-1}|\right) (1 \pm \delta_i) \leq \frac{\epsilon}{r} N e^{-(i-1)\epsilon} (1 + 2\delta) P \\ &\leq \frac{\epsilon}{r} N e^{-(i-1)\epsilon} (1 + 4\delta). \end{aligned}$$

# Proving the Theorem with the Nibble Lemma

$a > 0$ . Take  $\delta, \epsilon$  very small, and  $p$  s.t.  $e^{-\epsilon p} < \epsilon$ .

$$|V_i| \leq N e^{-i\epsilon}(1 + 2\delta) \quad \text{and} \quad |\mathcal{E}'_i| \leq \frac{\epsilon}{r} N e^{-(i-1)\epsilon}(1 + 4\delta).$$

Cover of size  $\sum_{i=1}^p |\mathcal{E}'_i| + |V_p|$ .

$$\begin{aligned} (1 + 4\delta) \frac{\epsilon}{r} N \sum_{i=0}^{p-1} e^{-i\epsilon} + |V_p| &\leq (1 + 4\delta) \frac{\epsilon N}{r} \frac{1}{1 - e^{-\epsilon}} + (1 + 2\delta) N e^{-\epsilon p} \\ &\leq \frac{N}{r} (1 + 4\delta) \left( \frac{1}{1 - e^{-\epsilon}} + r\epsilon \right) \\ &< (1 + a) \frac{N}{r}. \end{aligned}$$

# Proving the Nibble Lemma

$r > 2$  fixed.

For  $K \geq 1$ ,  $\epsilon > 0$  and  $\delta' > 0$ , there exists  $\delta > 0$  and  $D^*$  s.t. the following holds for all  $N \geq D \geq D^*$ .

Every  $r$ -uniform hypergraph  $H = (V, \mathcal{E})$  on  $N$  vertices such that

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contains a set  $\mathcal{E}'$  of hyperedges s.t.

- (iv)  $|\mathcal{E}'| = \frac{N}{r}\epsilon(1 \pm \delta')$ .
- (v)  $V' = V \setminus \bigcup_{e \in \mathcal{E}'} e$  has size  $Ne^{-\epsilon}(1 \pm \delta')$ .
- (vi) All vertices  $x$  of  $V'$  except at most  $\delta'|V'|$  the degree  $d'(x)$  of  $x$  in  $H[V']$  satisfies  $d'(x) = De^{-\epsilon(r-1)}(1 \pm \delta')$ .

# Proving the Nibble Lemma

$\mathcal{E}'$  a random subset of  $\mathcal{E}$  : every  $e \in \mathcal{E}$  is picked randomly, independently with probability  $p = \epsilon/D$ .

- With very high probability (say  $\geq 0.9$ ) (iv) holds.
- With very high probability (say  $\geq 0.9$ ) (v) holds.
- With very high probability (say  $\geq 0.9$ ) (vi) holds.

$\Rightarrow$  with positive probability ( $\geq 0.7$ ), (iv), (v) and (vi) hold.

With very high probability, (iv) holds

$$(iv) |\mathcal{E}'| = \frac{N}{r} \epsilon (1 \pm \delta').$$

(1): All vertices  $x$  except at most  $\delta N$  satisfy  $d(x) = (1 \pm \delta)D$ .

$$\implies \frac{1}{r}(1 - \delta)^2 DN \leq |\mathcal{E}| \leq \frac{1}{r}(1 + \delta)DN + \delta KDN$$

$$|\mathcal{E}| = (1 \pm \delta_1) \frac{DN}{r}.$$

$$\mathbf{E}(|\mathcal{E}'|) = |\mathcal{E}|/p = (1 \pm \delta_1) \frac{\epsilon N}{r}$$

**Need to prove** that  $|\mathcal{E}'|$  **is concentrated** around its expected value.

## 2nd Moment Method

**Variance** :  $\mathbf{Var}(X) = \mathbf{E} \left( (X - E(X))^2 \right)$ .

**Chebyshev Inequality** : For any  $\lambda$ ,

$$\Pr \left( |X - \mathbf{E}(X)| \geq \lambda \sqrt{\mathbf{Var}(X)} \right) \leq \frac{1}{\lambda^2} .$$

$$\Pr \left( X = \mathbf{E}(X) \pm \lambda \sqrt{\mathbf{Var}(X)} \right) \geq 1 - \frac{1}{\lambda^2} .$$

## Variance and covariance

Assume  $X = X_1 + \dots + X_m$ .

$$\mathbf{Var}(X) = \sum_{i=1}^m \mathbf{Var}(X_i) + \sum_{i \neq j} \mathbf{Cov}(X_i, X_j)$$

**Covariance** :  $\mathbf{Cov}(Y, Z) = \mathbf{E}(YZ) - \mathbf{E}(Y)\mathbf{E}(Z)$ .

If  $Y$  and  $Z$  are independent, then  $\mathbf{Cov}(Y, Z) = 0$ .

If  $X_i = 1$  with probability  $p_i$  and 0 otherwise, then

$$\mathbf{Var}(X_i) = p_i(1 - p_i) \leq p_i = \mathbf{E}(X_i)$$

$$\mathbf{Var}(X) \leq \mathbf{E}(X) + \sum_{i \neq j} \mathbf{Cov}(X_i, X_j)$$

With very high probability, (iv) holds

$$(iv) |\mathcal{E}'| = \frac{N}{r} \epsilon (1 \pm \delta').$$

(1): All vertices  $x$  except at most  $\delta N$  satisfy  $d(x) = (1 \pm \delta)D$ .

$$\implies \frac{1}{r}(1 - \delta)^2 DN \leq |\mathcal{E}| \leq \frac{1}{r}(1 + \delta)DN + \delta KDN$$

$$|\mathcal{E}| = (1 \pm \delta_1) \frac{DN}{r}.$$

$$\mathbf{E}(|\mathcal{E}'|) = |\mathcal{E}|/p = (1 \pm \delta_1) \frac{\epsilon N}{r}$$

$$\mathbf{Var}(|\mathcal{E}'|) = |\mathcal{E}|p(1 - p) \leq (1 \pm \delta_1) \frac{\epsilon N}{r}.$$

By Chebyshev Inequality,

$$\Pr \left( |\mathcal{E}'| = (1 \pm \delta_2) \frac{\epsilon N}{r} \right) > 0,9.$$



With very high probability, (v) holds

(v)  $V' = V \setminus \bigcup_{e \in \mathcal{E}'} e$  has size  $Ne^{-\epsilon}(1 \pm \delta')$ .

$I_x = 1$  if  $x \notin \bigcup_{e \in \mathcal{E}'} e$  and  $I_x = 0$  otherwise.

$$|V'| = \sum_{x \in V} I_x.$$

$x$  good if  $d(x) = (1 \pm \delta D)$ , bad otherwise.

$$\begin{aligned} x \text{ good : } \mathbf{E}(I_x) &= \mathbf{Pr}(I_x = 1) = (1 - p)^{d(x)} = \left(1 - \frac{\epsilon}{D}\right)^{(1 \pm \delta D)} \\ &= e^{-\epsilon}(1 \pm \delta_3). \end{aligned}$$

$x$  bad:  $0 \leq \mathbf{E}(I_x) \leq 1$ , but at most  $\delta N$  bad vertices.

**Linearity of the Expected Value :  $\mathbf{E}(|V'|) \leq Ne^{-\epsilon}(1 \pm \delta_4)$ .**

With very high probability, (v) holds

$$\begin{aligned}\mathbf{Var}(|V'|) &= \sum_{x \in V} \mathbf{Var}(I_x) + \sum_{x, y \in V, x \neq y} \mathbf{Cov}(I_x, I_y) \\ &\leq \mathbf{E}(|V'|) + \sum_{x, y \in V, x \neq y} \mathbf{Cov}(I_x, I_y)\end{aligned}$$

$$\begin{aligned}\mathbf{Cov}(I_x, I_y) &= \mathbf{E}(I_x I_y) - \mathbf{E}(I_x) \mathbf{E}(I_y) \\ &= (1 - p)^{d(x) + d(y) - d(x, y)} - (1 - p)^{d(x) + d(y)} \\ &\leq (1 - p)^{-d(x, y)} - 1 \leq \left(1 - \frac{\epsilon}{D}\right)^{-\delta(D)} - 1 \leq \delta_5.\end{aligned}$$

$$\mathbf{Var}(|V'|) \leq \mathbf{E}(|V'|) + \delta_5 N^2 \leq \delta_6 (\mathbf{E}(|V'|))^2 .$$

**Chebyshev** : With proba  $\geq 0,9$ ,

$$|V'| = (1 \pm \delta_7) \mathbf{E}(|V'|) = (1 \pm \delta_8) N e^{-\epsilon} .$$